

MATH 2073 -SOLUTION OF NONLINEAR EQUATIONS

LECTURE-3

Dejen Ketema

Department of Mathematics
Arba Minch University

<https://elearning.amu.edu.et/course/view.php?id=279>

Fall 2019



- The **root-finding problem** consists of the following: given a continuous function f , find the values of x that satisfy the equation

$$f(x) = 0. \quad (1)$$

- The solutions of this equation are called the **zeros** of f or the **roots** of the equation.
- In general, Eqn. (1) is impossible to solve exactly.
- Therefore, one must rely on some numerical methods for an approximate solution.
- The methods we will discuss in this chapter are iterative and consist basically of two types:
 - one in which the convergence is guaranteed and
 - the other in which the convergence depends on the initial guess.



STOPPING CRITERIA

The methods used for solving equations numerically can be divided into two groups: **bracketing methods** and **open methods**.

STOPPING CRITERIA

- Many ways to decide when to stop:
- Let c^* be the true (exact) solution such that $f(c^*) = 0$, and let m_k be a numerically approximated solution such that $f(m_k) = \epsilon$ (where ϵ is a small number).
- **Tolerance in the solution** A tolerance is the maximum amount by which the true solution can deviate from an approximate numerical solution.



WHEN DO WE STOP?

Let m_k is approximate solution of $f(x)$ at k^{th} iteration.

We can

- 1 keep going until successive iterates are close: $|m_k - m_{k-1}| < \epsilon$
- 2 close in relative terms: $\frac{|m_k - m_{k-1}|}{|m_k|} < \epsilon$
- 3 the function value is small enough: $|f(m_k)| < \epsilon$

No choice is perfect. In general, where no additional information about f is known, the second criterion is the preferred one (since it comes the closest to testing the relative error).



CONVERGENCE

Suppose an algorithm generates a sequence of approximations, c_n , which approaches a limit, c_* , or

$$\lim_{n \rightarrow \infty} c_n = c_*$$

How quickly does $c_n \rightarrow c_*$?

RATE OF CONVERGENCE

If a sequence c_1, c_2, \dots, c_n converges to a value c_* and if there exist real numbers $\lambda > 0$ and $\alpha \geq 1$ such that

$$\lim_{n \rightarrow \infty} \frac{|c_{n+1} - c_*|}{|c_n - c_*|^\alpha} = \lambda$$

then we say that α is the rate of **convergence of the sequence**.



THEOREM (THE INTERMEDIATE VALUE THEOREM)

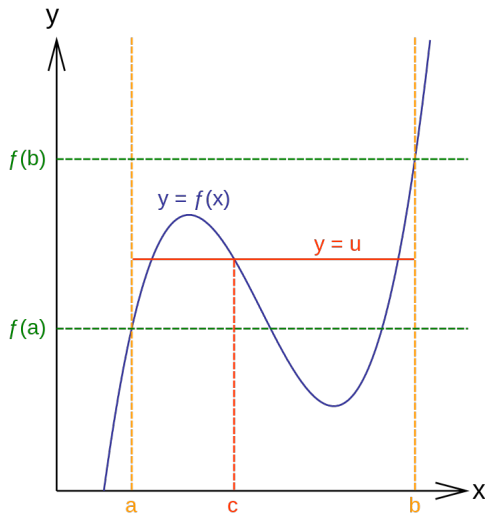
The intermediate value theorem states that if a continuous function, f , with an interval, $[a, b]$, as its domain, takes values $f(a)$ and $f(b)$ at each end of the interval, then it also takes any value between $f(a)$ and $f(b)$ at some point within the interval.

THEOREM (BOLZANO'S THEOREM)

If a continuous function has values of opposite sign inside an interval, then it has a root in that interval. Let $f : [a, b] \rightarrow \mathbb{R}$ be a continuous function such that $f(a) \cdot f(b) < 0$. Then there exist at least one point x in the open interval (a, b) such that $f(x) = 0$.



INTERMEDIATE VALUE THEOREM



BISECTION

The input for the method is a continuous function f , an interval $[a_0, b_0]$, and the function values $f(a_0)$ and $f(b_0)$. The function values are of opposite sign (there is at least one zero crossing within the interval).

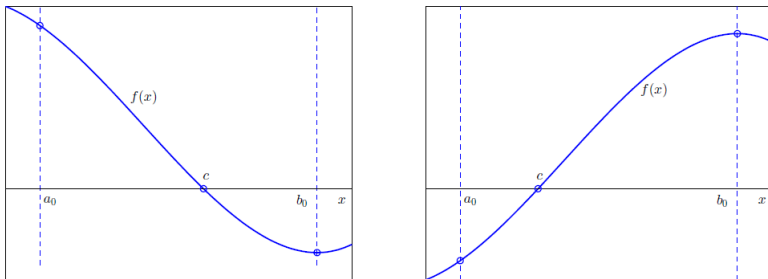


FIGURE: Bisection Method



THE BISECTION METHOD PROCEDURE IS:

- ➊ Choose a starting interval $[a_0, b_0]$ such that $f(a_0)f(b_0) < 0$.
- ➋ Compute $f(m_0)$ where $m_0 = (a_0 + b_0)/2$ is the midpoint.
- ➌ Determine the next subinterval $[a_1, b_1]$:
 - ➊ If $f(a_0)f(m_0) < 0$, then let $[a_1, b_1]$ be the next interval with $a_1 = a_0$ and $b_1 = m_0$.
 - ➋ If $f(b_0)f(m_0) < 0$, then let $[a_1, b_1]$ be the next interval with $a_1 = m_0$ and $b_1 = b_0$.
- ➍ Repeat (2) and (3) until the interval $[a_N, b_N]$ reaches some predetermined length.
- ➎ Constructs a sequence of intervals containing the root c :

$$(a_0, b_0) \supset (a_1, b_1) \supset \cdots \supset (a_{N-1}, b_{N-1}) \supset (a_N, b_N) \ni c$$

- ➏ After k steps

$$|b_k - a_k| = \frac{1}{2}|b_{k-1} - a_{k-1}| = \left(\frac{1}{2}\right)^k |b_0 - a_0|$$

- ➐ At step k , the midpoint $m_k = \frac{a_k + b_k}{2}$ is an estimate for the root c with

ABSOLUTE ERROR

The bisection method does not (in general) produce an exact solution of an equation $f(x) = 0$. However, we can give an estimate of the absolute error in the approximation.

THEOREM

Let $f(x)$ be a continuous function on $[a, b]$ such that $f(a)f(b) < 0$. After N iterations of the bisection method, let x_N be the midpoint in the N^{th} subinterval $[a_N, b_N]$

$$x_N = \frac{a_N + b_N}{2}$$

There exists an exact solution x_{true} of the equation $f(x) = 0$ in the subinterval $[a_N, b_N]$ and the absolute error is

$$|x_{\text{true}} - x_N| \leq \frac{b - a}{2^{N+1}}$$



Note that we can rearrange the error bound to see the minimum number of iterations required to guarantee absolute error less than a prescribed ϵ :

$$\frac{b-a}{2^{N+1}} < \epsilon$$

$$\frac{b-a}{\epsilon} < 2^{N+1}$$

$$\ln\left(\frac{b-a}{\epsilon}\right) < (N+1)\ln(2)$$

$$\frac{\ln\left(\frac{b-a}{\epsilon}\right)}{\ln(2)} - 1 < N$$



EXAMPLE: BISECTION METHOD

EXAMPLE

The bisection method applied to $f(x) = \left(\frac{x}{2}\right)^2 - \sin(x) = 0$ with $(a_0, b_0) = (1.5, 2.0)$, and $(f(a_0), f(b_0)) = (-0.4350, 0.0907)$ gives:

k	a_k	b_k	m_k	$f(m_k)$
0	1.5000	2.0000	1.7500	-0.2184
1	1.7500	2.0000	1.8750	-0.0752
2	1.8750	2.0000	1.9375	0.0050
3	1.8750	1.9375	1.9062	-0.0358
4	1.9062	1.9375	1.9219	-0.0156
5	1.9219	1.9375	1.9297	-0.0054
6	1.9297	1.9375	1.9336	-0.0002
7	1.9336	1.9375	1.9355	0.0024
8	1.9336	1.9355	1.9346	0.0011
9	1.9336	1.9346	1.9341	0.0004

ADVANTAGES, DISADVANTAGES

ADVANTAGES AND DISADVANTAGES

- 1 The method is guaranteed to converge
- 2 The error bound decreases by half with each iteration
- 3 The bisection method converges very slowly
- 4 The bisection method cannot detect multiple roots



EXAMPLE

Let's use our function with input parameters $f(x) = x^2 - x - 1$ and $N = 25$ iterations on $[1, 2]$ to approximate the **golden ratio**

$$\phi = \frac{1 + \sqrt{5}}{2}$$

The golden ratio is a root of the quadratic polynomial $x^2 - x - 1 = 0$.

$$1.618033990263939$$

The absolute error is guaranteed to be less than $(2 - 1)/(2^{26})$ which is

$$2^{(-26)} = 1.4901161193847656e^{-08}$$

